- What do we expect to find at all, if the sequence is long but finite?
- If we do not find an exact match of the word s, do we find at least a similar string?
- Do we find a high scoring local alignment of s and the database?
- How high scoring do we expect it to be?
- ullet This obviously depends on the word s.
- It also depends on the length of the random sequence.
- The longer we can search, the more chances we have for finding something.

### Some more notations

• To ensure local alignments, we now assume:

$$\sum_{i,j} \mu_i \, \mu_j \, S_{ij} < 0$$

and let

- Sequence1  $X_1, X_2, \ldots$ Sequence2  $Y_1, Y_2, \ldots$ be both i.i.d. sequences.
- We align the sequences (trivially and non optimally) by matching  $X_1$  with  $Y_1$ ,  $X_2$  with  $Y_2$ , and so on ...
- ullet Let  $Z_i=(X_i,Y_i)$  . We get a new i.i.d. sequence of real numbers  $S(Z_1),S(Z_2),\ldots$ , where  $S(Z_i)$  is the score of matching  $X_i$  and  $Y_I$ .
- ullet Let  $\mu$  be the distribution of  $Z_i$ , we have

$$E_{\mu}[S(Z_i)] = m_0 < 0$$

# Large Deviations

By the Law of large numbers we have:

$$\frac{1}{n} \sum_{k=1}^{n} S(Z_k) \longrightarrow m_0$$

For  $m>m_0$  let

$$A_n(m) := \left\{ \frac{1}{n} \sum_{k=1}^n S(Z_k) > m \right\}.$$

Again by the law of large numbers we have

$$\lim_{n} P\left[A_n(m)\right] = 0.$$

Events of this type are called large deviations.

# High scoring alignments

- Let us assume that the first n
  positions of the two random sequences
  are very similar...
  - ... more precisely let us assume they have a positive score.
- This event is a large deviation and hence, it becomes more and more unlikely the longer the sequences are.
- With which rate does the probability of this event converge to zero ?

### Cramer's theorem

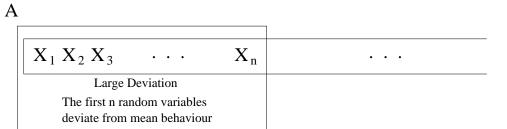
• Theorem: (Cramér 37)

$$\lim_{n \to \infty} \frac{1}{n} \log \left( P \left[ A_n(m) \right] \right) = -I(m, S, \mu) \quad a.s. .$$

- Probabilities of large deviations decrease exponentially:  $P\left[A_n(m)\right] \approx \exp\left(-n\,I(m,S,\mu)\right).$
- $\bullet$  There is an explicit formula for  $I(m,S,\mu) \ \dots \ \ \text{see notes on the web.}$

### Local Scores

- The above alignments start at the first position of the sequences.
- However, local alignments can start every where inside the sequences.
- For our problem, we are not interested in the probability that the start segment of the alignment is high scoring, but in the probability that there is at least one high scoring segment inside the alignment.
- We search the optimal scoring segment in the fixed alignment.
- The corresponding score is called the maximal local score of the sequence  $S(Z_1), S(Z_2), \ldots$



### The Problem:

• For an i.i.d. sequence of real valued random variables we have the following problem:

What is the probability that the sequence contains a segment with an average well above the expectation?

### Rare Blocks

- Consider a long segment of the sequence  $Z_1, Z_2, \ldots$  given by a start point k and an end point l > k.
- If k and l are selected randomly and if the length of the segment l-k is long enough we expect by the law of large numbers:

$$\frac{1}{l-k} \sum_{i=k}^{l} S(Z_i) \approx m_0.$$

• Suppose the average score inside the segment is  $m>m_0$ . This can be described by the event

$$B_{k,l}(m) = \left\{ \frac{1}{l-k} \sum_{i=k}^{l} S(Z_i) > m \right\}.$$

Such events are called rare blocks. For fixed (k,l) Cramér's theorem yields

$$P\left[B_{k,l}(m)\right] \approx e^{-|l-k|I(m,S,\mu)}$$
.

## Long rare blocks

- We will often observe short segments in the fixed alignment with a positive score. However, long stretches of similarity are rare.
- We are interested in the length  $R_n$  of the longest rare block inside a sequences of length n.
- Long sequences provide more possible starting points for a rare block than short ones.
- Therefore,  $R_n$  grows with the length of the i.i.d. sequence under consideration.
- What is the relationship between  $R_n$  and the length of  $S(Z_1), S(Z_2), \ldots$ ?

# Erdoes-Renyi Law

• The answer is given by  $\begin{array}{lll} \textbf{Theorem:} & (\text{Erd\"{o}s R\'{e}nyi 1970}) \text{ Let} \\ S(Z_1), S(Z_2), \dots \text{ be i.i.d.} & \text{with common} \\ \text{mean } m_0. & \text{For } m > m_o \text{ let } R_n \text{ denote} \\ \text{the length of the longest m-block in} \\ (S(Z_1), \dots, S(Z_n)). & \text{We have} \end{array}$ 

$$\lim_{n o \infty} rac{R_n}{\log(n)} = rac{1}{I(m,S,\mu)}$$
 a. s.

- The key observation is that the maximal local score is significantly higher than the score of the sequence beginning at position 1.
   Consequently, it can be interpreted as a large deviation in the context of Cramer's theorem.
- See the notes for more details. The notes also come back to the problem of waiting times for rare blocks.

 Rare blocks grow on a logarithmic rate with the length of the underlying sequence.

### Head Runs

- ullet We toss a fair coin n times.
- The possible outcomes are Head (H) and Tail (T).

### 

- How long is the longest head run?
- Erdoes Renyi: It depends on the length n of the sequence. The length of the longest head run is proportional to log(n).

### Exact matches

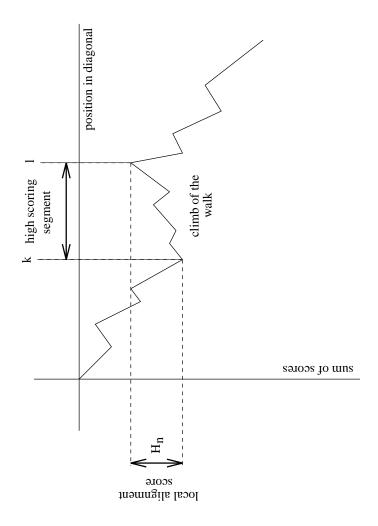
- We have a fixed alignment of two random sequences:
- Match is head, mismatch is tail: TATAGCATCATTACT CATGGATCCATTCGA THHTHTTTHHHHTTT
- The length of the longest match is proportional to the log of the length of the alignment.

# Climbs of a random walk with negative drift

• Consider the partial sums

$$W_n = \sum_{i=1}^n S(Z_i).$$

• Since the expectation of the  $C_i$  is negative, the sequence of partial sums  $W_1,W_2,\ldots$ , yields a random walk with a negative drift.



• With probability one we have

$$\lim_{n\to\infty} W_n = -\infty,$$

and therefore for  $l\gg k$  we have  $W_l-W_k<0$  with high probability.

- ullet Now suppose we have a segment with  $H=W_l-W_k=h_0>0$ . Note that this segment corresponds to a gap-free local alignment with score  $h_0$ .
- The setup is very similar as in the Erdös-Rènyi law: ''typical'' behavior is a decreasing path, but since the walk is random there are also some short excursions of the walk, in which it is climbing against its drift.
- The main difference in the concept is, that the deviation is not determined by the length of the climb, but by its height.

# Large deviations for local scores

The two large deviation problems of high alignment scores and rare blocks are closely related, and it is not surprising that they have similar answers.

Theorem:(Dembo and Karlin 1991) Let  $W_1, W_2, \ldots$  be a random walk with discrete i.i.d. increments  $S(Z_i)$  that have a negative expectation. If  $H_n$  denotes the height of the highest climb in  $W_1, \ldots, W_n$ , i.e.

$$H_n = \max_{l,k} W_l - W_k,$$

we have

$$\lim_{n o\infty}rac{H_n}{\log(n)}= heta$$
 a. s. ,

where  $\theta$  is the unique positive solution of the equation

$$E\left[\exp(S(Z_1)/\theta)\right] = 1.$$

# Optimal alignment

 Up to know, we have fixed the alignment

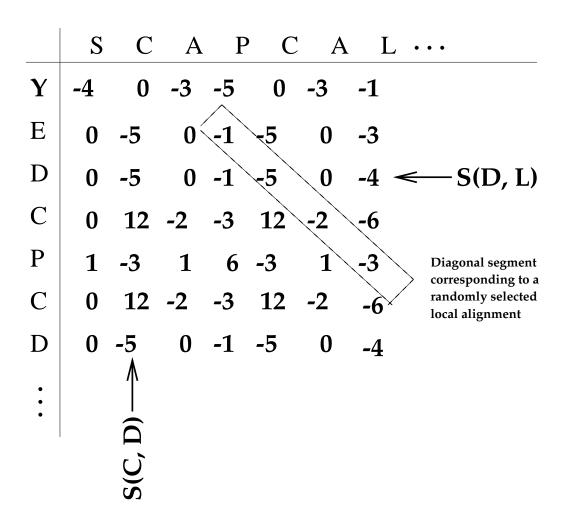
```
X1 X2 X3 X4 ...
```

Y1 Y2 Y3 Y4 ...

- That is not how real alignments are made.
- Let us now proceed to local alignments without gaps.
- We choose among all possible pairs of segments the one that optimizes the score.

# Calculate a table C with

$$C_{i,j} = S(X_i, Y_j).$$



- There is a one to one correspondence of gap-free local alignments and segments of diagonals in this table.
- To each pair of equally long segments in the sequences there exists a corresponding segment of a diagonal, such that the sum over all entries in this diagonal segment yields the corresponding local alignment score.
- ullet If the local alignment starts at positions  $X_i$  and  $Y_j$  and has length L its score is given by the sum over the diagonal segment

$$(C_{i,j}, C_{i+1,j+1}, \ldots, C_{i+L-1,j+L-1})$$

- Now, we are talking optimal local alignments without gaps.
- Rare blocks correspond to long alignments with a score above the average and climbs of a random walk correspond to high scoring alignments
- The main difference is that there are nm possible starting points for a rare block or a climb of a random walk.
- The diagonals are dependent ... a technical problem that can be overcome.

Dembo Karlin (1994)

- Erdoes Renyi: The length of the longest common word of two sequences of length n and m grows proportional to  $\log(n\,m)$
- Dembo Karlin: The optimal local ungapped alignment score grows proportional to  $log(n\,m)$ .
- The rates are different.

### Some more Head Runs

- Match is head, mismatch is tail: TATAGCATCATTACT CATGGATCCATTCGA THHTHTTTHHHHTTT
- Let us fix a head run length t such that head runs of this length are rare.
- Let W(t) be the number of head runs that are longer or equal to t.
- ullet W(t) is a count of rare events ...
- ... does it follow a Poisson distribution?

# Clumping Head runs

t=5

Clump of head runs that have length 5

# Declumping Head runs

t=5

# Poisson approximation

- The number of long head runs is not Poisson ... because of clumps
- We declump
- The number of clumps is Poisson distributed
- Not trivial because clumps are not independent
- Chen-Stein Theorem ... see Waterman's book for more details.

- Probability of a head is p and of a tail (1-p).
- ullet The length of the sequence is L
- Fix a minimal length t for a head run and a position n in the sequence.
- What is the probability P that a clump starts at n?
  - 1. n = 1:  $P = p^t$
  - 2.  $1 < n \le L t$ :  $P = (1 p)p^t$
  - 3. n > L t P = 0

• The expected number of clumps in the sequence

$$\lambda(t) := E[W(t)] = p^t \{ (n-t)(1-p) + 1 \}.$$

 $\bullet$  Let  $Z(\lambda)$  be a Poisson variable with intensity  $\lambda$ 

$$P[Z(\lambda) = k] = \frac{e^{-\lambda} \lambda^k}{k!},$$

$$E[Z(\lambda)] = \lambda$$

• Approximate W(t) by  $Z(\lambda(t))$ .

### Extreme value distribution

 If the longest head run is shorter than t, the count of head runs longer or equal to t is zero.

$$\{R_n < t\} = \{W(t) = 0\}$$

and therefore

$$P[R_n \ge t] \approx 1 - e^{-\lambda(t)}$$
.

• ... and therefore

$$P[R_n \ge t] \approx 1 - e^{-\lambda_n(t)}.$$

• For  $\theta = -1/\log(p)$ , large t, and  $(n-t)/n \approx 1$  we have

$$P[R_n \ge t] \approx 1 - \exp\left(-p^t \{(n-t)(1-p) + 1\}\right)$$
$$\approx 1 - \exp\left(-\gamma n e^{-\frac{t}{\theta}}\right), \tag{1}$$

- Let  $H_{nm}$  be the optimal ungapped local alignment score of two i.i.d. sequences with length n and m.
- Along the same lines as above, we can get

$$P[H_{nm} > t] \approx 1 - \exp(-\gamma \, m \, n \, e^{-\frac{t}{\theta}})$$

ullet Of course the parameters  $\gamma$  and  $\theta$  are different.

ullet A standard extreme value distributed random variable G is defined by

$$P[G < t] = \exp(-e^{-t}).$$

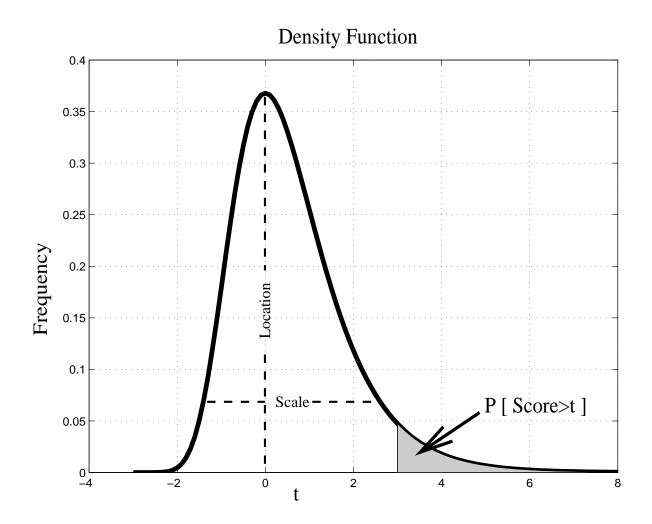
 Shifting and rescaling yields the two parameter family of variables

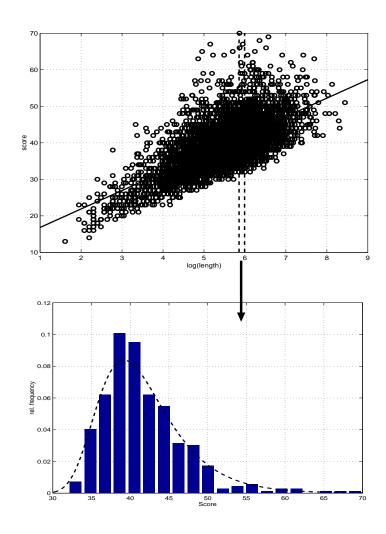
$$H = \theta G + \xi$$

We call  $\xi$  the location and  $\theta$  the scale of H.

 With this notation we have the regression equation

$$H_{nm} \sim \alpha + \theta \log(n m) + \theta G$$
.





- Note, that the location of the score distribution, depends on the length of sequences
- In the comparisons of a database search, the length of the query is fixed but the lengths of the database entries vary.
- Long sequences are more likely to produce a high score, than short ones
- Long sequences tend to show up as false positives in the twilight zone.
- We should compensate for this effect.

# Length corrections

Rank sequences according to the p-values

$$P[H_{nm} > t] \approx 1 - \exp(-\gamma \, m \, n \, e^{-\frac{t}{\theta}})$$

2. ... or rank them according to length adjusted scores

$$A = H_{nm} - log(nm)\theta.$$